

Exercises for the Lecture Logics  
Sheet 5

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Delivery until 25. Mai 2011 10:00 Uhr

**Exercise 1:** [Tableaux, tutorial]

On sheet 1 a diet plan was described by the following formula:

$$A \equiv (\neg B \rightarrow F) \wedge (((B \wedge F) \rightarrow \neg E) \wedge ((E \vee \neg B) \rightarrow \neg F))$$

Construct a complete tableau for  $A$ . What properties of  $A$  can be derived using the tableau? Find a disjunctive normal form for  $A$  using the tableau.

**Exercise 2:** [Tableaux-consequence, tutorial]

Prove:

1.  $(A \wedge \neg B) \vdash_{\tau} \neg((\neg A) \wedge (\neg B))$
2.  $(A \wedge (A \rightarrow B)) \vdash_{\tau} B$
3.  $A \rightarrow (B \rightarrow C) \vdash_{\tau} (A \rightarrow B) \rightarrow (A \rightarrow C)$

**Exercise 3:** [compactness theorem, 6P]

1. Prove that for a set of formulas  $\Sigma$  and a formula  $A$  there is a closed tableau for  $\Sigma \cup \{\neg A\}$  iff there is a closed tableau for  $\Gamma \cup \{\neg A\}$  for every finite subset  $\Gamma \subset \Sigma$ .  
Hint: You can use König's Lemma, which states that a tree is finite iff every branch is finite.
2. Prove that the compactness theorem holds for tableaux, i.e.

$$\Sigma \vdash_{\tau} A \text{ iff there is a finite subset } \Sigma_0 \subseteq \Sigma \text{ with } \Sigma_0 \vdash_{\tau} A.$$

**Exercise 4:** [Tableaux-consequence, 5P]

Prove:

1.  $\{p, p \vee q, p \rightarrow s, r \rightarrow q\} \vdash_{\tau} q \rightarrow p$
2.  $\{p, p \vee q, p \rightarrow s, r \rightarrow q\} \vdash_{\tau} s$
3.  $\vdash_{\tau} (\neg(p \rightarrow q) \rightarrow (q \rightarrow p))$
4.  $F \vdash_{\tau} q \rightarrow p \wedge (\neg(s \wedge \neg(s \vee ((q \wedge r) \rightarrow p))))$
5.  $\neg((A \rightarrow (A \vee C)) \wedge D) \vdash_{\tau} (C \rightarrow B) \vee \neg D$

**Exercise 5:** [DNF from tableaux, 2P]

Using the tableaux-method, find a DNF for the following formulas:

1.  $(p \rightarrow \neg(\neg q \rightarrow r)) \rightarrow (q \vee r)$
2.  $(q \rightarrow p) \rightarrow \neg(r \rightarrow q)$

**Exercise 6:** [Tableaux with equivalence, 4P]

$\alpha$ - and  $\beta$ -formulas have so far only been defined for  $\{\neg, \wedge, \vee, \rightarrow\}$ , but  $\leftrightarrow$  has been omitted. Is  $A \equiv B \leftrightarrow C$  an  $\alpha$ - or a  $\beta$ -formula and which components does it consist of?

**Exercise 7:** [limitations of the tableau method, 5P]

Can the following statements be proved using the tableaux method? Give short explanations.

1.  $\{p, q, r, s\} \models t$
2.  $\{p, q, r, s\} \not\models \neg(q \rightarrow s)$
3.  $F \models \neg(p \rightarrow (q \leftrightarrow r) \wedge \neg r) \rightarrow (s \vee \neg p)$
4.  $\Sigma := \{p_i \wedge \neg p_{i+1} \mid i \in \mathbb{N}\}$  is unsatisfiable.
5.  $\Sigma := \{p_i \wedge \neg p_{i+1} \mid i \in 2\mathbb{N}\}$  is satisfiable.

Keep your arguments as general as possible. I.e. if you can explain why there is a tableaux-proof without actually writing it down, then do not write it down.

**Exercise 8:** [Tableaux-consequence, tutorial]

Prove without using the soundness and completeness of the tableaux-method:

If  $\Sigma \vdash_{\mathcal{T}} p \wedge q$  then  $\Sigma \vdash_{\mathcal{T}} p \vee q$  holds as well.

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