

## Chapter 6

# Verifying Functions

# Motivation

## Verification

Verifying properties of functions is a fundamental task in SE. Hence it is an aspect of theorem proving. In particular, functions definitions allow to express recursive algorithms. Our focus here is on the definition of:

- ▶ termination/well-definedness properties
- ▶ functional properties, i.e., properties relating input parameters to the result (PR-properties).
- ▶ Example: A compiler can be considered as a partial function.

## In general:

- ▶ specification = model + properties
- ▶ or
- ▶ specification = model\_1 + model\_2 + relationship

# Conceptual aspects

Here: specification = function definition + PR-properties

Verify:

- ▶ well-definedness of function by:
  - ▶ often structural induction according to parameter types
  - ▶ more general: well-founded ordering on parameter space “show that parameters become smaller”
- ▶ PR-properties:
  - ▶ often structural induction according to parameter types
  - ▶ in general, proof technique depends on properties

# Discussion

## Verification

- ▶ works for the full parameter space (in contrast to testing)
- ▶ checks for consistency of models and properties
  - ▶ models may not reflect what programmer had in mind
  - ▶ properties may not reflect what programmer had in mind
  - ▶ proofs can have errors
- ▶ uses redundancy to find errors
- ▶ helps to improve the descriptions

# Discussion (cont.)

## Formal verification

- ▶ avoids misunderstanding
- ▶ allows using tools
- ▶ avoids errors in proofs
- ▶  $\rightsquigarrow$  Isabelle and others



# Case study: greatest common divisor

see Gcd.thy

# Case study: Quicksort

## Assumptions

Given:

```
datatype mapping = lt | ge

fun eval :: "mapping => universe => universe => bool"
  where
    "eval ge xa ya = not(eval lt ya xa)" |
    "[|eval lt ya xa|] ==> eval lt xa ya = False"
```

Modeling in Isabelle using type classes!



# Case study: Quicksort

Shallow embedding of the algorithm:

## Case study: Quicksort (cont.)

```

fun qsplrit  ::
  "mapping => universe => universe list => universe
    list"
  where

  "qsplrit xf xa Nil      = Nil" |
  "qsplrit xf xa (ya#x) =
    (if eval xf ya xa then ya#qsplrit xf xa x
      else qsplrit xf xa x)"

fun qsort :: "universe list => universe list" where
  "qsort Nil      = Nil" |
  "qsort (p # l) =
    qsort (qsplrit lt p l) @ p # qsort (qsplrit ge p l)"

```

# Properties to prove

## Well-definedness/termination of qsort (1) and qsplit (2)

```
primrec counts :: "'a list => 'a => nat" where
  "counts [] x      = 0" |
  "counts (y#yl) x = counts yl x +(if x=y then 1 else 0)"
  "
```

## lemma qsort\_counts(3): “counts xl = counts (qsort xl)”

```
fun qsorted :: "universe list => bool" where
  "qsorted [] = True" |
  "qsorted [x] = True" |
  "qsorted (a#b#l) = (ge b a \and qsorted (b#l))"
```

## lemma qsort\_sort\_prop(4): “qsorted (qsort xl)”

# Verification of the properties

Ad 1: qsplit is primitive recursive

Ad 2: Idea: length of parameter decreases

```
Auxiliary lemma qsplit_length:  
  "length (qsplit f p l) <= length l"
```

↪ Proof termination with “length” as measure

## Verification of the properties (cont.)

Auxiliary lemma `counts_concat`:

```
"counts (l @ m) x = (counts l x) + (counts m x)"
```

Auxiliary theorem `qsplit_lt_ge_count [iff]`:

```
"count (qsplit lt p l) x + count (qsplit ge p l) x =  
  count l x"
```

Prove lemma “`qsort_counts`” by induction

# Property 4

## Order lifting to lists

```
primrec qall :: "mapping => universe => universe list
  => bool" where
  "qall f p [] = True"
| "qall f p (h # t) = (f h p \and qall f p t)"
```

## Property 4 (cont.)

### Auxiliary Properties

```
theorem qsplit_splits:
  "qall f p (qsplit f p l)"

lemma qall_concat :
  "qall f p (a @ b) = (qall f p a \and qall f p b)"

theorem qsplit_qall :
  "qall f p l ==> qall f p (qsplit g q l)"

theorem qsort_qall :
  "qall f p l ==> qall f p (qsort l)"
```

## prop(4): “sorted (qsort xl)”

### Auxiliary lemmas

```
lemma qsorted_append :  
  "[| qsorted l; qall ge p l|] ==> qsorted (p # l)"
```

```
theorem qsorted_concat :  
  "[| qsorted a; qsorted b; qall lt p a; qall ge p b  
   |] ==> qsorted (a @ p # b)"
```

»> Generic.QSort.thy