

## Exercises to the Lecture FSVT

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sheet 8

**Exercise 1:** [Confluence and termination of rule sets over ground terms]Let  $R = \{(l_k, r_k) | k = 1, \dots, n\}$  be a finite rule set over ground terms. Prove:

1. If there is an infinite chain, then there is a rule  $(l, r) \in R$  with an infinite chain from  $r$ .
2. If there is an infinite chain, then there is a  $j$  with  $1 \leq j \leq n$  and a ground term  $t$ , such that  $l_j \xrightarrow{\pm} t$  and  $l_j$  is a subterm of  $t$ .
3. Termination of  $R$  is decidable. (Termination is often denounced as 'Kettenbedingung' in german literature.)
4. Develop sufficient conditions for local confluence.

**Exercise 2:** [Knuth-Bendix-ordering]Let  $\varphi : F \cup V \rightarrow \mathbb{N}$  be a weight function with

$$\varphi(x) = \alpha > 0 \quad \text{for all } x \in V \quad (1)$$

$$\varphi(f) \geq \alpha \quad \text{if } f \text{ 0-ary} \quad (2)$$

$$\varphi(f) > 0 \quad \text{if } f \text{ 1-ary} \quad (3)$$

$$\varphi(f) \geq 0 \quad \text{else} \quad (4)$$

Extend  $\varphi$  to  $\varphi : \text{Term}(F, V) \rightarrow \mathbb{N}$  by

$$\varphi(f(t_1, \dots, t_n)) = \varphi(f) + \sum_{i=1, \dots, n} \varphi(t_i)$$

Define  $s > t$  iff.  $\varphi(s) > \varphi(t)$  and  $|s|_x \geq |t|_x$  for all  $x \in V$ . Then  $>$  is called a Knuth-Bendix-ordering. Prove for any Knuth-Bendix-ordering  $>$ :

1.  $>$  is strict part of a wellfounded partial ordering
2.  $>$  is compatible with substitution
3.  $>$  is compatible with term replacement

**Exercise 3:**

Let

$$R_1 = \{F(0, 1, x) \rightarrow F(x, x, x)\}$$

$$R_2 = \{G(x, y) \rightarrow x, G(x, y) \rightarrow y\}.$$

1. Prove:  $R_1$  and  $R_2$  are terminating.
2. Prove or disprove: The rule set  $R_1 \cup R_2$  is terminating.

**Exercise 4:** [Example confluence and critical pairs]

Consider the rule system  $R : h(x, f(x)) \rightarrow c, h(x, x) \rightarrow b, k(x) \rightarrow x, g(a) \rightarrow f(g(k(a)))$ .

1. Prove: There are no critical pairs of  $R$ .
2. Prove:  $R$  is not confluent.
3. Why is there no contradiction?

**Exercise 5:** [Local coherence and critical pairs]

Prove: Let  $\text{CP}(R, G)$  be defined as the set of critical pairs regarding  $R$  and the set of equations  $G$  oriented in both ways. If  $R$  is left-linear, then the following statements are equivalent.

1.  $\rightarrow_R$  is locally coherent modulo  $\sim$ .
2. For every critical pair  $(t_1, t_2) \in \text{CP}(R, G)$  holds  $t_1 \downarrow_{\sim} t_2$ .

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